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Multivariate Tropical Regression and Piecewise-Linear Surface Fitting

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Optimal Regression for Fitting Euclidean vs Tropical Lines

Problem: Fit a curve to data (x_i, y_i) , i = 1, ..., m

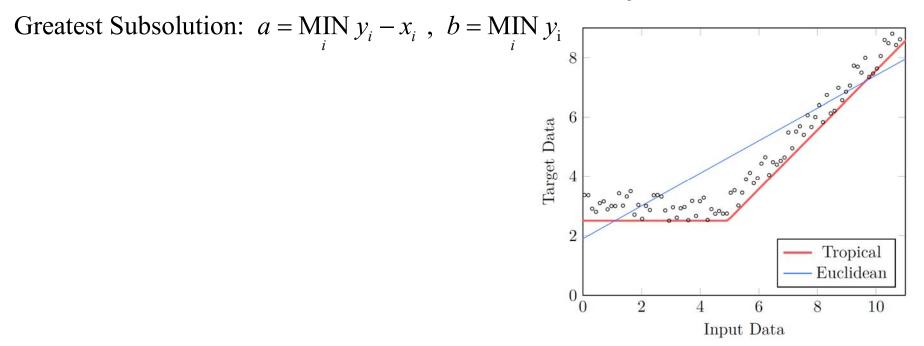
Euclidean:

Fit a straight line y = ax + b by minimizing ℓ_2 -norm of error:

$$a = \frac{\sum x_i y_i - (\sum x_i) (\sum y_i) / m}{\sum (x_i)^2 - (\sum x_i)^2 / m} , \ b = \frac{1}{m} \sum_i y_i - a x_i$$

Tropical:

Fit a tropical line $y = \max(a + x, b)$ by minimizing some ℓ_p -norm of error:



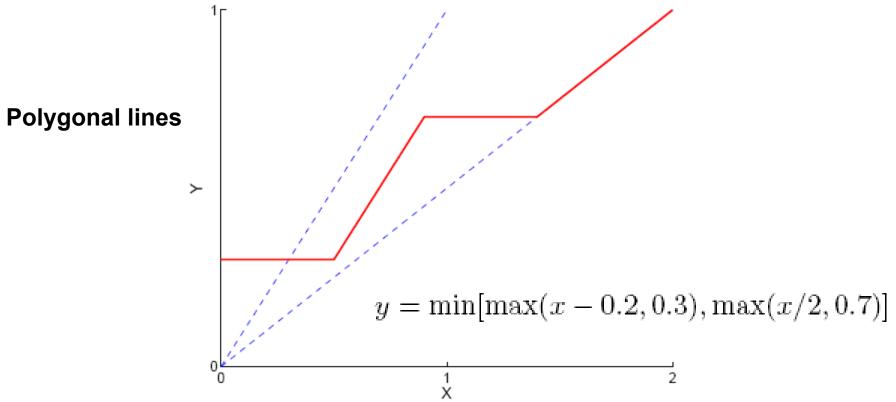
Outline

1. Elements of Tropical Geometry

- "a marriage between algebraic geometry and polyhedral geometry" [Maclagan & Sturmfels 2015]
- Tropical semirings: Max-plus & Min-plus Arithmetic
- Tropical Polynomials
- Geometrical objects: tropical lines, tropical polynomial curves/surfaces, tropical half-spaces & polyhedra
- 2. Elements of Max-plus Algebra for Vectors/Matrices:
 - Max-plus Matrix algebra
 - Max-plus systems (Nonlinear Control) and
 - Signal Processing: Max-plus convolutions (weighted dilations/erosions)
- **3. Optimization and Tropical Regression:**
 - Optimal solutions of max-plus matrix equations
 - Tropical Regression: fitting tropical polynomials to data
 - Algorithm and Complexity

What does TROPICAL mean?

- The adjective "tropical" was coined by French mathematicians Dominique Perrin and Jean-Eric Pin, to honor their Brazilian colleague Imre Simon, a pioneer of min-plus algebra as applied to finite automata in computer science.
- Tropical (Τροπικός in Greek) comes from the greek word «Τροπή» which means "turning" or "changing the way/direction".



Tropical Semirings

Scalar Arithmetic Rings

Integer/Real Addition-Multiplication Ring: $(\mathbb{R}, +, \times), (\mathbb{Z}, +, \times)$

Tropical Semirings

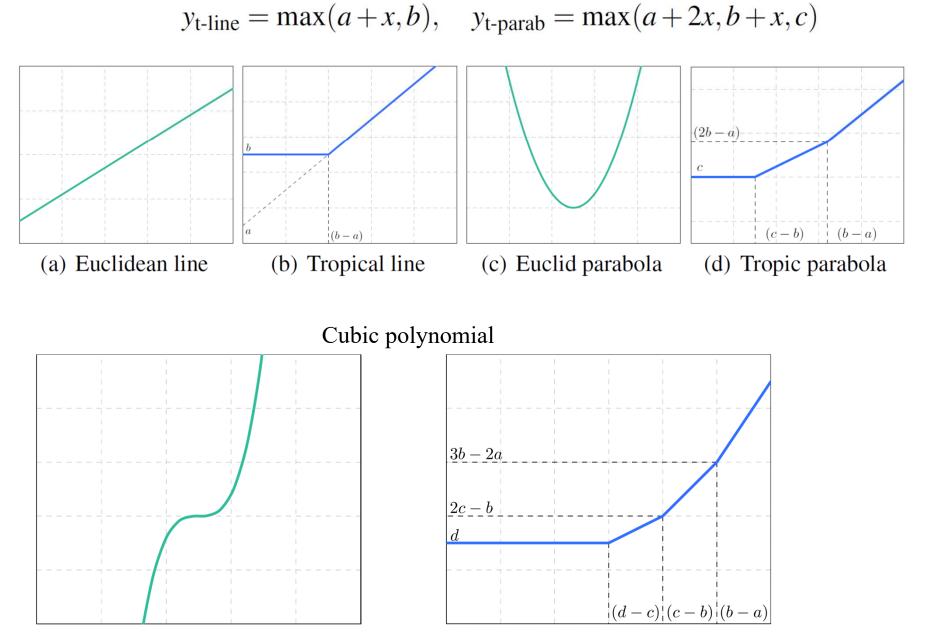
$$\mathbb{R}_{\max} = \mathbb{R} \cup \{-\infty\}, \ \mathbb{R}_{\min} = \mathbb{R} \cup \{+\infty\}$$
$$\lor = \max, \ \land = \min$$
Max-plus semiring: $(\mathbb{R}_{\max}, \lor, +)$ Min-plus semiring: $(\mathbb{R}_{\min}, \land, +)$

Log-Sum-Exp (LSE) approximation

(Maslov "Dequantization" in idempotent mathematics [Maslov 1987, Litvinov 2007])

$$\lim_{T \downarrow 0} T \cdot \log(e^{a/T} + e^{b/T}) = \max(a, b)$$
$$\lim_{T \downarrow 0} (-T) \log(e^{-a/T} + e^{-b/T}) = \min(a, b)$$

Graphs of Max-plus Tropical 1D Polynomials

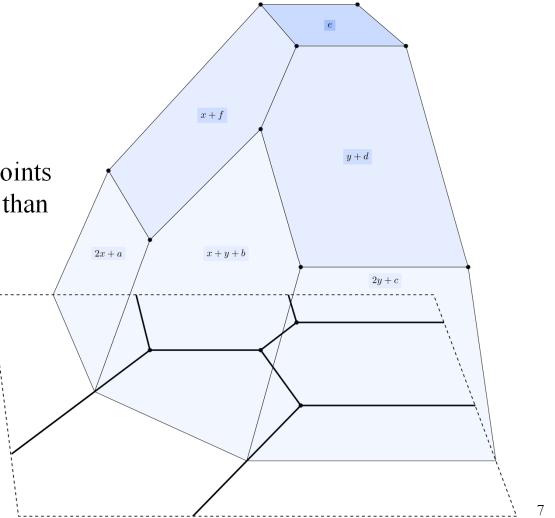


Graph and Trop Curve of a tropical "Conic" polynomial

Tropical Polynomial of degree 2 in two variables

classical: $ax^{2} + bxy + cy^{2} + dy + e + fx''$ tropical: $p(x, y) = \min(a + 2x, b + x + y, c + 2y, d + y, e, f + x)$

Graph of p(x,y)and its **Tropical Curve** = set of (x,y) points where the min is attained by more than one terms.



Obtain Tropical Polynomials via Dequantization

Classic polynomial:
$$f(\mathbf{u}) = \sum_{k=1}^{K} c_k u_1^{a_{k1}} u_2^{a_{k2}} \cdots u_n^{a_{kn}}, \quad \mathbf{u} = (u_1, u_2, \dots, u_n)$$

Posynomial if $c_k > 0$, $\mathbf{a}_k = (a_{k1}, \dots, a_{kn}) \in \mathbb{R}^n$, $\mathbf{u} > 0$;

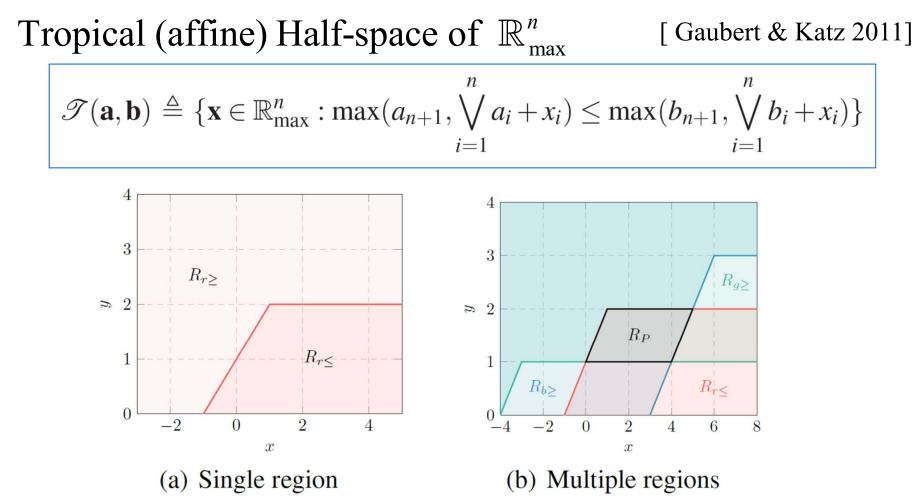
Log-Sum-Exp (Viro's "logarithmic paper" [Viro 2001]):

$$\mathbf{x} = \log(\mathbf{u}), \quad b_k = \log(c_k)$$
$$\lim_{T \downarrow 0} T \cdot \log f(e^{\mathbf{x}/T}) = \lim_{T \downarrow 0} T \cdot \sum_{k=1}^K \exp(\langle \mathbf{a}_k, \mathbf{x} / T \rangle + b_k / T) \rightarrow$$

Tropical (max-plus) **Polynomial** = Piecewise-Linear Function

$$p(\mathbf{x}) = \bigwedge_{k=1}^{K} \left\{ \langle \mathbf{a}_{k}, \mathbf{x} \rangle + b_{k} \right\} = \bigwedge_{k=1}^{K} \left\{ a_{k1} x_{1} + \dots + a_{kn} x_{n} + b_{k} \right\}$$

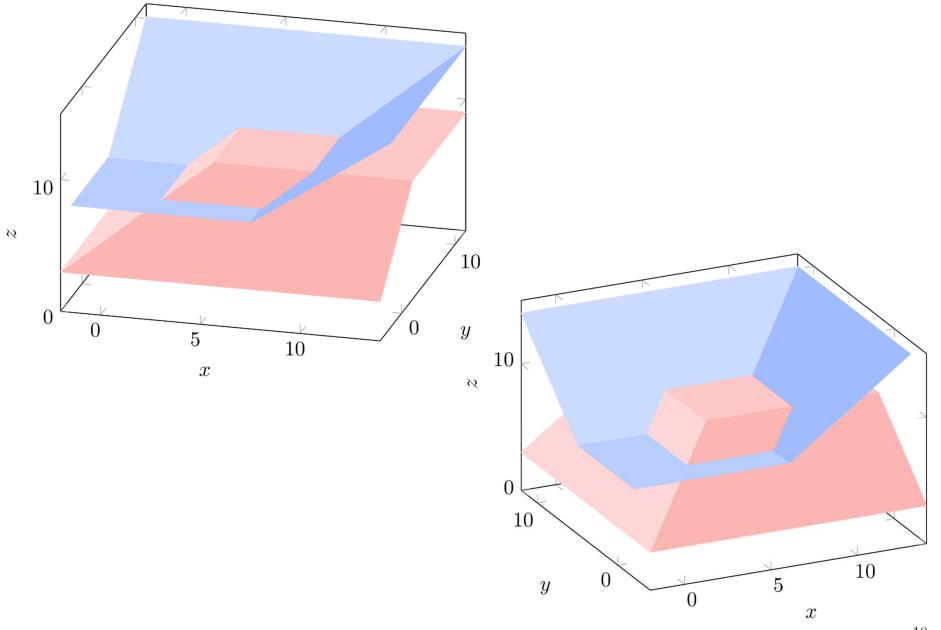
Tropical Half-spaces and Polytopes in 2D



The region separating boundaries are tropical lines (or hyper-planes).

Tropical **Polyhedra** are formed from finite intersections of tropical half-spaces. **Polytopes** are compact polyhedra.

Tropical Halfspaces and Polyhedra in 3D



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Maxplus Matrix Algebra

• vector/matrix 'addition' = pointwise max

$$\mathbf{x} \lor \mathbf{y} = [x_1 \lor y_1, \dots, x_n \lor y_n]^T$$

$$\mathbf{A} \lor \mathbf{B} = [a_{ij} \lor b_{ij}]$$

• vector/matrix 'dual addition' = pointwise min

$$\mathbf{x} \wedge \mathbf{y} = [x_1 \wedge y_1, \dots, x_n \wedge y_n]^T \mathbf{A} \wedge \mathbf{B} = [a_{ij} \wedge b_{ij}]$$

• vector/matrix 'multiplication by scalar'

$$c + \mathbf{x} = [c + x_1, \dots, c + x_n]^T$$

$$c + \mathbf{A} = [c + a_{ij}]$$

• (max, +) 'matrix multiplication'

$$[\mathbf{A} \boxplus \mathbf{B}]_{ij} = \bigvee_{k=1}^{n} a_{ik} + b_{kj}$$

• (min, +) 'matrix dual multiplication'

$$[\mathbf{A} \boxplus' \mathbf{B}]_{ij} = \bigwedge_{k=1}^{n} a_{ik} + b_{kj}$$

Linear versus Max-Plus Systems

- State-space representation: linear vs. max-plus
 - $\begin{aligned} x\left(k\right) &= Ax\left(k-1\right) + Bu\left(k\right) \\ y\left(k\right) &= Cx\left(k\right) + Du\left(k\right) \end{aligned} \qquad \begin{aligned} x\left(k\right) &= A \boxplus x\left(k-1\right) \lor B \boxplus u\left(k\right) \\ y\left(k\right) &= C \boxplus x\left(k\right) \lor D \boxplus u\left(k\right) \end{aligned}$
- Matrix products
 - Linear: $[AB]_{ij} = \sum_{k=1}^{n} a_{ik} b_{kj}$
 - Max-plus: $\overline{[A \boxplus B]_{ij}} = \bigvee_{k=1}^{n} a_{ik} + b_{kj}$
- Example

$$\begin{bmatrix} 4 & -1 \\ 2 & -\infty \end{bmatrix} \boxplus \begin{bmatrix} x \\ y \end{bmatrix} = \begin{bmatrix} 3 \\ 1 \end{bmatrix}, \left\{ \begin{array}{c} \max(x+4, y-1) = 3 \\ x+2 = 1 \end{array} \right\} \Longrightarrow \begin{array}{c} x = -1 \\ y \le 4 \end{array}$$

• What can we model with max-plus systems?

Discrete Event Systems/Control, Scheduling, Shortest Paths on Graphs, Dynamic Programming, WFSTs for Speech recognition, Morphological Filters for Image Processing

Ref: [P. Maragos, "Dynamical Systems on Weighted Lattices: General Theory", Math. Control, Signals and Systems, 2017.]

Morphological Operators on Lattices

 $(\leq = partial ordering, V = supremum, \Lambda = infimum)$

- ψ is increasing iff $f \leq g \Rightarrow \psi(f) \leq \psi(g)$.
- δ is **dilation** iff $\delta(\vee_i f_i) = \vee_i \delta(f_i)$.
- ε is erosion iff $\varepsilon(\wedge_i f_i) = \wedge_i \varepsilon(f_i)$.
- α is **opening** iff increasing and antiextensive $(\alpha(f) \le f)$, and idempotent $(\alpha = \alpha^2)$.
- β is **closing** iff increasing and extensive $(\beta(f) \ge f)$, and idempotent $(\beta = \beta^2)$.
- (ε, δ) is adjunction iff $\delta(g) \le f \Leftrightarrow g \le \varepsilon(f)$.

(Galois connection) Residuation pair ("Tropical Adjoints")

Then: ε is erosion, δ is dilation,

 $\delta \varepsilon$ is opening (projection), $\varepsilon \delta$ is closing (projection).

[Serra 1988; Heijmans & Ronse 1990]

Solve Max-plus Equations

- Problems:
 - (1) Exact problem: Solve $\delta_A(\mathbf{x}) = \mathbf{A} \boxplus \mathbf{x} = \mathbf{b}, \quad \mathbf{A} \in \overline{\mathbb{R}}^{m \times n}, \ \mathbf{b} \in \overline{\mathbb{R}}^m$
 - (2) Approximate Constrained: Min $\|\mathbf{A} \boxplus \mathbf{x} \mathbf{b}\|_{p=1...\infty}$ s.t. $\mathbf{A} \boxplus \mathbf{x} \leq \mathbf{b}$
- Theorem: (a) The greatest (sub)solution of (1) and unique solution of (2) is

$$\hat{\mathbf{x}} = \varepsilon_A(\mathbf{b}) = \mathbf{A}^* \boxplus' \mathbf{b} = [\bigwedge_i b_i - a_{ij}], \quad \mathbf{A}^* \triangleq -\mathbf{A}^T$$

and yields the **Greatest Lower Estimate (GLE)** of data **b**:

$$\delta_A(\varepsilon_A(\mathbf{b})) = \mathbf{A} \boxplus (\mathbf{A}^* \boxplus' \mathbf{b}) \le \mathbf{b}$$

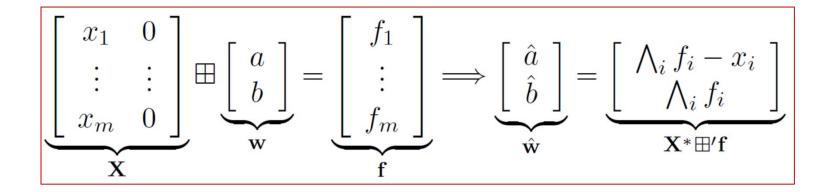
(b) Min Max Absolute Error (MMAE) unconstrained unique solution:

$$\tilde{\mathbf{x}} = \hat{\mathbf{x}} + \mu, \quad \mu = \|\mathbf{A} \boxplus \hat{\mathbf{x}} - \mathbf{b}\|_{\infty}/2$$

- Geometry: Operators δ, ε are vector dilation and erosion, and the GLE $\mathbf{b} \mapsto \delta \varepsilon(\mathbf{b})$ is an opening (lattice projection).
- **Complexity**: O(mn)

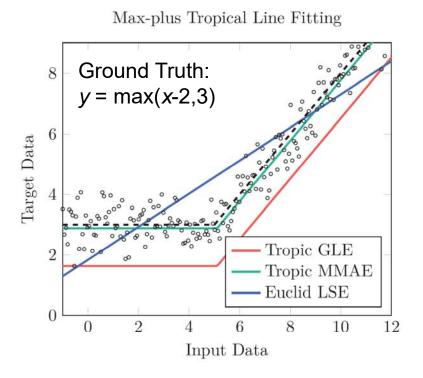
Optimally Fitting Tropical Lines to Data

Problem: Fit a tropical line $y = \max(a + x, b)$ to noisy data (x_i, f_i) , i = 1, ..., m, where $f_i = y_i$ +error by minimizing $\ell_{1,...,\infty}$ norm of error: **Greatest Subsolution (GLE)**: $\hat{w} = (\hat{a}, \hat{b})$, $\hat{a} = \min_i f_i - x_i$, $\hat{b} = \min_i f_i$ Min Max Abs. Error (MMAE) Solution: $\tilde{w} = \hat{w} + \mu$, $\mu = ||$ GLE error $||_{\infty}/2$

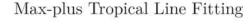


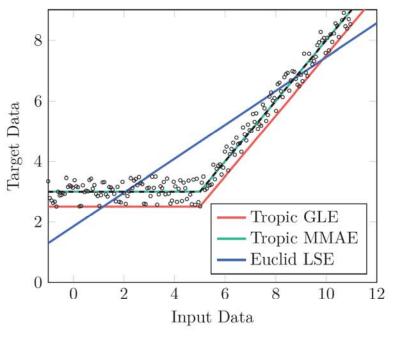
Numerical Examples of Optimally Fitting Tropical Lines to Data

Problem: Fit a tropical line $y = \max(a + x, b)$ to noisy data (x_i, f_i) , i = 1, ..., m = 200, where $f_i = y_i$ +error by minimizing $\ell_{1,...,\infty}$ of error: Greatest Subsolution (GLE): $\hat{w} = (\hat{a}, \hat{b})$, $\hat{a} = \min_i f_i - x_i$, $\hat{b} = \min_i f_i$ Min Max Abs. Error (MMAE) Solution: $\tilde{w} = \hat{w} + \mu$, $\mu = ||$ GLE error $||_{\infty} / 2$



(a) T-line with Gaussian Noise



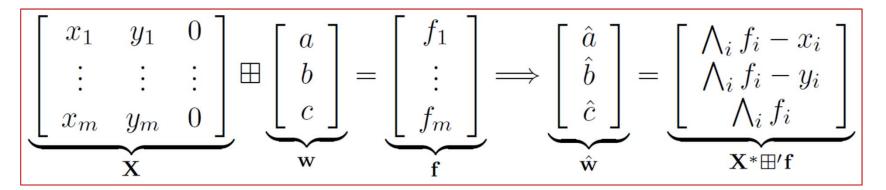


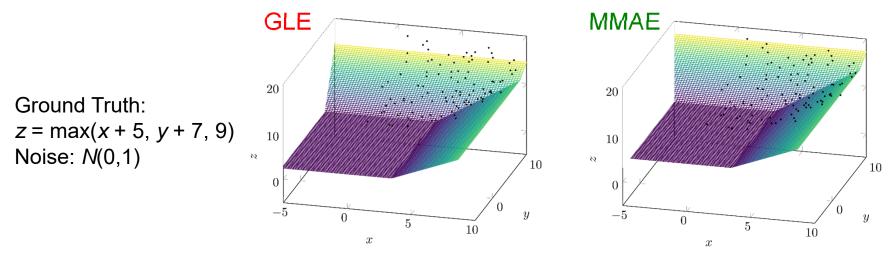
(b) T-line with Uniform Noise

Optimal Fitting Max-Plus Tropical Planes to Data

Problem: Fit a tropical plane $z = \max(a + x, b + y, c)$ to noisy data (x_i, y_i, f_i) , where $f_i = z_i$ +error, i = 1, ..., m = 100, by minimizing $\ell_{1,...,\infty}$ norm of error: **Greatest Subsolution (GLE)**: $\hat{w} = (\hat{a}, \hat{b}, \hat{c})$

Min Max Abs. Error (MMAE) Solution: $\tilde{w} = \hat{w} + \mu$, $\mu = || \text{ GLE error } ||_{\infty} / 2$





Optimal Fitting 2D Higher-degree Tropical Polynomials to Data

Data (noisy paraboloid): 3D tuples $(x_i, y_i, f_i) \in \mathbb{R}^3$ $f_i = x_i^2 + y_i^2 + \varepsilon_i$, $(x_i, y_i) \sim \text{Unif}[-1,1]$ $\varepsilon_i \sim \mathcal{N}(0, 0.25^2)$

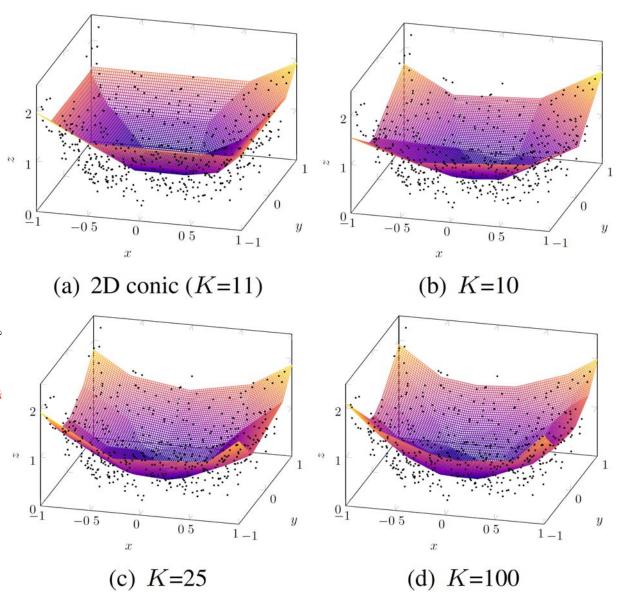
Model:

Fit K-rank 2D trop. polynomial

 $p(x, y) = \underset{k=1}{\overset{K}{\operatorname{AX}}} \left\{ a_{k} x + b_{k} y + c_{k} \right\}$

by minimizing error $||f_i - p(x_i, y_i)||_{\circ}$ Estimation algorithm:

K – means on data gradients $\rightarrow a_k, b_k$ solve max-plus eqns $\rightarrow c_k$



Computational Complexity of Tropical Regression

Data size: *m* data in \mathbb{R}^{n+1} **Model**: MAX of *K* hyperplanes $\mathbf{a}_k^T \cdot \mathbf{x} + b_k$, k = 1, ..., K, $\mathbf{x} \in \mathbb{R}^n$ # of **Parameters**: K(n+1)

Traditional Least-Squares Estimator: $O((n+1)^3 m^3)$ (Quadratic Programming with Constraints)

Iterative *K*-means Partition and LSE: $O((n+1)^2 mi)$ [Magnani & Boyd 2009], i = # iterations until convergence

Convex Adaptive Partitioning: $O(n(n+1)^2 m \log(m) \log(\log(m)))$ [Hannah & Dunson 2011]

Our algorithm: O(Kmni)i = # K-means iterations if $iK \ll m \implies O(mn)$

Conclusions

- Tropical Geometry, max-plus matrix algebra, and morphological signal operators share a common idempotent semiring arithmetic.
- Introduced Tropical Polynomials for multidimensional data fitting using Piecewise-Linear Functions.
- Developed algorithm of low-complexity (~linear) for tropical regression based on optimal solutions of systems of max-plus equations.
- Future work: extensions to more general regression functions using max-* algebra on weighted lattices.

References

- P. Maragos and E. Theodosis, "*Tropical Geometry and Piecewise-Linear* Approximation of Curves and Surfaces on Weighted Lattices", arXiv, 2019.
- P. Maragos, "*Dynamical Systems on Weighted Lattices: General Theory*", Math. Control, Signals and Systems, 2017.

Thank you for your attention!

We wish everyone courage and health during the COVID19 pandemic.

For more information, demos, and current results:

http://cvsp.cs.ntua.gr and http://robotics.ntua.gr